### Markov Chain Monte Carlo: Applications to MIMO detection and channel equalization

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# **Background – MCMC methods**

- Markov Chain Monte Carlo (MCMC) is a statistical method to generate random samples from arbitrary distributions.
- Previous work on MCMC methods for signal processing and communication (Doucet-Wang'05):
  - Bit-counting: use MCMC simulations to determine the frequency over which a bit occurs.
  - Many samples are needed.
  - Requires a burning period to allow Markov chain to converge.
- Our proposed MCMC methods:
  - Do not use bit-counting
  - No burning period is needed
  - Require very few samples even for large systems

# **Background– comparisons of suboptimal detectors**

- Linear detectors: zero-forcing (ZF), minimum mean square error (MMSE): low complexity, limited performance
- List Sphere Decoding (LSD) detectors (Hochwald-Brink'03), Soft-in Softout (SISO)-LSD (Vikalo-Hassibi-Kailath'04):
  - $\blacklozenge$  Use tree search to find a sample set  $\mathcal B$  containing likely samples.
  - Excellent performance at high SNR
  - Variable complexity
  - Exponential complexity at low SNR
- MCMC detectors:
  - $\blacklozenge$  Use Gibbs sampler to find  $\mathcal B$
  - Constant complexity
  - Excellent performance at low SNR
  - High SNR problem

# System diagram– Joint MIMO detection and channel decoding



• Channel Model:

$$\mathbf{y} = \sqrt{\frac{\rho}{t}} \mathbf{H} \mathbf{d} + \mathbf{n}, \tag{1}$$

- t transmit and r receive antenna
- $\rho$ : SNR per receive antenna.
- **H**: *r* by *t* channel fading matrix, i.i.d. complex Gaussian CN(0, 1);
- $\mathbf{d} = (d_1, \cdots, d_t)^T$ : transmitted signal vector
- $\mathbf{y} = (y_1, \cdots, y_r)^T$ : received signal vector
- **n**: i.i.d. CN(0, 1) distributed entries;
- Assumes coherent detection: receiver knows H perfectly.
- Noncoherent MCMC detector (Chen-Peng'09)

#### **Optimal detector**

- $\mathbf{d} = (d_1, \cdots, d_t)^T \leftrightarrow \mathbf{b} = (b_0, b_1, \cdots, b_{K-1})^T$ , where  $K = tM_c$ .
- $\lambda = (\lambda_0, \dots, \lambda_{K-1})^T$ ,  $\lambda_i$  is LLR of the *i*-th bit provided by channel decoder.
- Given y and  $\lambda$ , the LLR of bit  $b_k$  is

$$\gamma_{k} = \ln \frac{P(b_{k} = 1 | \mathbf{y}, \lambda)}{P(b_{k} = -1 | \mathbf{y}, \lambda)} = \ln \frac{\sum_{\overline{\mathbf{b}}_{k}} P(b_{k} = 1, \overline{\mathbf{b}}_{k} | \mathbf{y}, \lambda)}{\sum_{\overline{\mathbf{b}}_{k}} P(b_{k} = -1, \overline{\mathbf{b}}_{k} | \mathbf{y}, \lambda)}$$
(2)

where  $\overline{\mathbf{b}}_k = (b_0, \cdots , b_{k-1}, b_{k+1}, \cdots, b_{K-1}); b_j \in \{1, -1\}.$ 

• Optimal detector has exponential complexity  $2^{K}$ .

• Gibbs sampler – an example: t = 2, QPSK

|           |                                | TX1   |       | TX2   |       |                       |
|-----------|--------------------------------|-------|-------|-------|-------|-----------------------|
|           |                                | $b_0$ | $b_1$ | $b_2$ | $b_3$ |                       |
|           | $\mathbf{b}^{(0)} \rightarrow$ | 1     | -1    | 1     | 1     | random initialization |
| 1st       |                                | -1    | -1    | 1     | 1     | update $b_0$          |
| iteration |                                | -1    | 1     | 1     | 1     | update $b_1$          |
|           |                                | -1    | 1     | 1     | 1     | update $b_2$          |
|           | $\mathbf{b}^{(1)} \rightarrow$ | -1    | 1     | 1     | -1    | update b <sub>3</sub> |
| 2nd       |                                | -1    | 1     | 1     | -1    | update $b_0$          |
| iteration |                                | -1    | -1    | 1     | -1    | update $b_1$          |
|           |                                | -1    | -1    | 1     | -1    | update $b_2$          |
|           | $\mathbf{b}^{(2)} \rightarrow$ | -1    | -1    | 1     | 1     | update $b_3$          |

- Run *Q* Gibbs sampler in parallel, with *I* iterations each.
- Sample set  $\mathcal{B} = {\mathbf{b}^{(1)}, \mathbf{b}^{(2)}, \cdots, \mathbf{b}^{(QI)}}$ . Remove redundant samples.

• Max-Log:

$$\begin{aligned} \mathsf{LLR}_{k} &= \ln \frac{P(b_{k} = 1 | \mathbf{y}, \lambda)}{P(b_{k} = -1 | \mathbf{y}, \lambda)} \\ &\approx \max_{\{\mathbf{b}: \ \mathbf{b} \in \mathcal{B}_{+1}^{k}\}} \Big\{ - \left\| \mathbf{y} - \sqrt{\frac{\rho}{t}} \mathbf{Hd}(\mathbf{b}) \right\|^{2} + \frac{1}{2} \lambda^{T} \mathbf{b} \Big\} - \max_{\{\mathbf{b}: \ \mathbf{b} \in \mathcal{B}_{-1}^{k}\}} \Big\{ - \left\| \mathbf{y} - \sqrt{\frac{\rho}{t}} \mathbf{Hd}(\mathbf{b}) \right\|^{2} + \frac{1}{2} \lambda^{T} \mathbf{b} \Big\}. \end{aligned}$$

• Expanded set:

$$\mathcal{B} = \boxed{\begin{array}{rrrrr} -1 & 1 & 1 & -1 \\ -1 & -1 & 1 & 1 \end{array}}$$

then the expanded set for bit  $b_0$  is

$$\mathcal{B}^{0} = \frac{ \begin{array}{c|cccc} -1 & 1 & 1 & -1 \\ \hline 1 & 1 & 1 & -1 \\ \hline -1 & -1 & 1 & 1 \\ \hline 1 & -1 & 1 & 1 \end{array} = \mathcal{B}^{0}_{-1} \cup \mathcal{B}^{0}_{+1}$$

• Log-MAP-table-tb (Chen-Peng-Ashikhmin-Farhang'08):

$$LLR_{k} \approx \ln \sum_{\{\mathbf{b}: \mathbf{b} \in \mathcal{B}_{+1}^{k}\}} \exp\left\{-\left\|\mathbf{y} - \sqrt{\frac{\rho}{t}}\mathbf{Hd}(\mathbf{b})\right\|^{2} + \frac{1}{2}\lambda^{T}\mathbf{b}\right\}$$

$$-\ln \sum_{\{\mathbf{b}: \mathbf{b} \in \mathcal{B}_{-1}^{k}\}} \exp\left\{-\left\|\mathbf{y} - \sqrt{\frac{\rho}{t}}\mathbf{Hd}(\mathbf{b})\right\|^{2} + \frac{1}{2}\lambda^{T}\mathbf{b}\right\}.$$
(3)

To further reduce complexity,

$$\ln(e^{\delta_1} + e^{\delta_2}) = \max(\delta_1, \delta_2) + \ln(1 + e^{-|\delta_2 - \delta_1|})$$
  
= max(\delta\_1, \delta\_2) + f\_c(|\delta\_1 - \delta\_2|), (4)

#### Compared to Max-Log

- better performance
- less samples

# **b-MCMC MIMO detector – Simulation results**



Figure 1: Performance of turbo and LDPC coded TX8 64QAM systems [Chen-Peng-Ashikhmin-Farhang'08].

- Compare Max-Log, Log-MAP-tb b-MCMC, Max-Log LSD (Hochwald-Brink'03), Max-Log-SISO LSD (Vikalo-Hassibi-Kailath'04) and their Log-MAP versions
- Log-MAP-tb 10x10 MCMC performs the best in performance and complexity
- Within 1.8 dB of capacity at 24 bits/channel use.

To alleviate high SNR problems:

- Assume a larger noise variance than actual noise variance (Farhang-Zhu-Shi'06).
- Initialize one of the Gibbs sampler using ZF or MMSE solutions (Mao-Amini-Farhang'07).
- Constrained MCMC (Akoum-Peng-Chen-Farhang'09):

$$\mathsf{LLR}_{k} = \max_{\{\mathbf{b}: b_{k}=1\}} \left\{ -\left\| \mathbf{y} - \sqrt{\frac{\rho}{t}} \mathbf{Hd}(\mathbf{b}) \right\|^{2} \right\} - \max_{\{\mathbf{b}: b_{k}=-1\}} \left\{ -\left\| \mathbf{y} - \sqrt{\frac{\rho}{t}} \mathbf{Hd}(\mathbf{b}) \right\|^{2} \right\},$$
(5)

• First find ML solution  $b_{ML}^{(0)}$ .

• If  $b_{ML,k}^{(0)} = 1$ ,  $b_{ML}^{(0)}$  achieves the first maximum.

- N-ML solution: the vector that attains the second maximum.
- run constrained MCMC for each bit k to obtain good approximations of N-ML.

# **MCMC** equalizers for frequency selective channels



Channel Model:

$$y_n = \sum_{l=0}^{L} h_l x_{n-l} + z_n, \quad n = 0, 1, \cdots, N + L - 1,$$
 (6)

- *L*: channel memory
- $h_l$ : channel gain of *l*-th tap
- { $x_0, \cdots, x_{N-1}$ }: transmitted symbols
- $\{y_0, \cdots, y_{N-1}\}$ : received signals
- { $z_n$ }: i.i.d. channel noise  $CN(0, N_0/2)$ .

Group MCMC (g-MCMC) (Peng-Chen-Farhang'09):

- Inside Gibbs sampler, update  $G_{max}$  symbols at a time.
- g-MCMC performs better than b-MCMC for channels with strong ISI.

An example: Assume L = 2,  $G_{max} = 2$ , QPSK modulation.

|           |                                | <i>x</i> <sub>0</sub> | $x_1$ | <i>x</i> <sub>2</sub> | <i>x</i> <sub>3</sub> | <i>x</i> <sub>4</sub> | <i>x</i> <sub>5</sub> | <i>x</i> <sub>6</sub> | <i>X</i> <sub>7</sub> |                     |
|-----------|--------------------------------|-----------------------|-------|-----------------------|-----------------------|-----------------------|-----------------------|-----------------------|-----------------------|---------------------|
| 1st       | $\mathbf{x}^{(0)} \rightarrow$ | 1                     | 3     | 2                     | 2                     | 0                     | 3                     | 1                     | 0                     | initial vector      |
| iteration |                                | 3                     | 0     | 2                     | 2                     | 0                     | 3                     | 1                     | 0                     | update $(x_0, x_1)$ |
|           |                                | 3                     | 0     | 1                     | 2                     | 0                     | 3                     | 1                     | 0                     | update $(x_2, x_3)$ |
|           |                                | 3                     | 0     | 1                     | 2                     | 2                     | 3                     | 1                     | 0                     | update $(x_4, x_5)$ |
|           | $\mathbf{x}^{(1)} \rightarrow$ | 3                     | 0     | 1                     | 2                     | 2                     | 3                     | 1                     | 3                     | update $(x_6, x_7)$ |
| 2nd       |                                | 0                     | 0     | 1                     | 2                     | 2                     | 3                     | 1                     | 3                     | update $(x_0)$      |
| iteration |                                | 0                     | 1     | 1                     | 2                     | 2                     | 3                     | 1                     | 3                     | update $(x_1, x_2)$ |
|           |                                | 0                     | 1     | 1                     | 0                     | 3                     | 3                     | 1                     | 3                     | update $(x_3, x_4)$ |
|           |                                | 0                     | 1     | 1                     | 0                     | 3                     | 2                     | 0                     | 3                     | update $(x_5, x_6)$ |
|           | $\mathbf{x}^{(2)} \rightarrow$ | 0                     | 1     | 1                     | 0                     | 3                     | 2                     | 0                     | 0                     | update $(x_7)$      |

Compute LLR for bit  $b_k$ :

$$\gamma_{k} = \ln \frac{\sum_{i_{1}:i_{2} \in \mathcal{B}_{i_{1}:i_{2}}^{k,+1}} p(\mathbf{y}_{i:i+L} | \mathbf{b}_{i_{1}:i_{2}}) \prod_{l=i_{1}}^{i_{2}} P(b_{l})}{\sum_{\mathbf{b}_{i_{1}:i_{2}} \in \mathcal{B}_{i_{1}:i_{2}}^{k,-1}} p(\mathbf{y}_{i:i+L} | \mathbf{b}_{i_{1}:i_{2}}) \prod_{l=i_{1}}^{i_{2}} P(b_{l})}$$
(7)

- Bit k is mapped to symbol x<sub>i</sub>
- $\mathbf{y}_{i:i+L}$ : received signals that are affected by  $b_k$ .
- $\mathbf{y}_{i:i+L}$  depends only on bits  $\{b_l, i_1 = M_b(i-L) \le l \le M_b(i+L) = i_2\}$ .

A channel with strong ISI  $h_1[n] = 0.227\delta[n] + 0.46\delta[n-1] + 0.688\delta[n-2] + 0.46\delta[n-3] + 0.227\delta[n-4]$ 



Figure 2: Performance comparisons of MAP, MMSE, g-MCMC equalizers for a strong ISI channel [Chen-Peng-Farhang'09].

- g-MCMC significantly outperforms MMSE equalizer (Tuchler-Singer-Koetter'02)
- b-MCMC does not work for such channel with strong ISI.

# **Conclusions and future work**

- MCMC detectors are capable of achieving excellent performance at low complexity for both MIMO channels and frequency selective channels.
- Amicable for hardware implementation (Laraway-Farhang'09)
- MCMC equalizers allow for parallel implementation (Peng-Chen-Farhang'09)
- Ongoing research:
  - MCMC detectors for continuously time-varying channels and channels with imperfect channel state information.
  - Applications of MCMC equalizers to underwater acoustic channels.
  - Theoretical analysis of MCMC techniques.

# Thank You!