

## An Example: MIPS

From the Harris/Weste book  
Based on the MIPS-like processor from  
the Hennessy/Patterson book

## Instruction Set

Table 1.7 MIPS instruction set (subset supported)

| Instruction       | Function   | Encoding | op     | funct  |
|-------------------|--|----------|--------|--------|
| add \$1, \$2, \$3 | addition: \$1 → \$2 + \$3                                | R        | 000000 | 100000 |
| sub \$1, \$2, \$3 | subtraction: \$1 → \$2 - \$3                             | R        | 000000 | 100010 |
| and \$1, \$2, \$3 | bitwise and: \$1 → \$2 and \$3                           | R        | 000000 | 100100 |
| or \$1, \$2, \$3  | bitwise or: \$1 → \$2 or \$3                             | R        | 000000 | 100101 |
| slt \$1, \$2, \$3 | set less than: \$1 → 1 if \$2 < \$3<br>\$1 → 0 otherwise | R        | 000000 | 101010 |
| addi \$1, \$2,    | add immediate: \$1 → \$2 + imm                           | I        | 001000 | n/a    |
| beq \$1, \$2, imm | branch if equal: PC → PC + imm <sup>a</sup>              | I        | 000100 | n/a    |
| j destination     | jump: PC <sub>destination</sub> <sup>a</sup>             | J        | 000010 | n/a    |
| lb \$1, imm(\$2)  | load byte: \$1 → mem[\$2 + imm]                          | I        | 100000 | n/a    |
| sb \$1, imm(\$2)  | store byte: mem[\$2 + imm] → \$1                         | I        | 110000 | n/a    |

## Fibonacci (C)

$$f_0 = 1; f_{-1} = -1$$

$$f_n = f_{n-1} + f_{n-2}$$

$$f = 1, 1, 2, 3, 5, 8, 13, \dots$$

```
int fib(void)
{
    int n = 8;          /* compute nth Fibonacci number */
    int f1 = 1, f2 = -1; /* last two Fibonacci numbers */

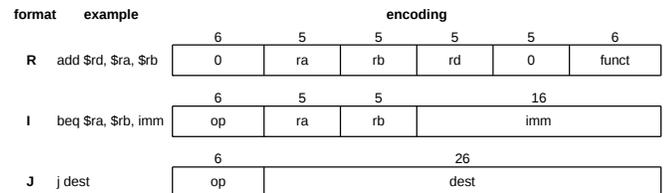
    while (n != 0) {   /* count down to n = 0 */
        f1 = f1 + f2;
        f2 = f1 - f2;
        n = n - 1;
    }
    return f1;
}
```

## MIPS Architecture

- Example: subset of MIPS processor architecture
  - Drawn from Patterson & Hennessy
- MIPS is a 32-bit architecture with 32 registers
  - Consider 8-bit subset using 8-bit datapath
  - Only implement 8 registers (\$0 - \$7)
  - \$0 hardwired to 00000000
  - 8-bit program counter

## Instruction Encoding

- 32-bit instruction encoding
  - Requires four cycles to fetch on 8-bit datapath



## Fibonacci (Assembly)

- 1<sup>st</sup> statement: int n = 8;
- How do we translate this to assembly?
  - Decide which register should hold its value
  - load an immediate value into that register
  - But, there's no "load immediate" instruction...
  - But, there is an addi instruction, and there's a convenient register that's always pinned to 0
- addi \$3, \$0, 8 ; load 0+8 into register 3

# Fibonacci (Assembly)

```
# fib.asm
# Register usage: $3: n $4: f $5: fp
# return value written to address 255
fib: addi $3, $0, 8 # initialize n=8
     addi $4, $0, 0 # initialize f = 0
     addi $5, $0, 1 # initialize fp = 1
loop: beq $3, $0, end # Done with loop if n == 0
     add $4, $4, $5 # f = f + fp
     sub $5, $4, $5 # fp = f - fp
     beq $3, $0, end
     addi $3, $3, -1 # n = n - 1
     j loop # while loop
end: sb $4, 255($0) # store result in address 255
```

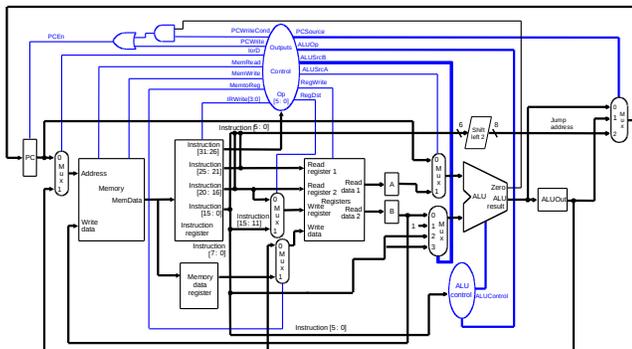
# Fibonacci (Binary)

## Machine language program

| Instruction       | Binary Encoding                       | Hexadecimal Encoding |
|-------------------|---------------------------------------|----------------------|
| addi \$3, \$0, 8  | 001000 00000 00011 0000000000001000   | 20030008             |
| addi \$4, \$0, 1  | 001000 00000 00100 0000000000000001   | 20040001             |
| addi \$5, \$0, -1 | 001000 00000 00101 1111111111111111   | 2005ffff             |
| beq \$3, \$0, end | 000100 00011 00000 0000000000000101   | 10600005             |
| add \$4, \$4, \$5 | 000000 00100 00101 00100 00000 100000 | 00852020             |
| sub \$5, \$4, \$5 | 000000 00100 00101 00101 00000 100010 | 00852822             |
| addi \$3, \$3, -1 | 001000 00011 00011 1111111111111111   | 2063ffff             |
| j loop            | 000010 0000000000000000000000000011   | 08000003             |
| sb \$4, 255(\$0)  | 110000 00000 00100 0000000011111111   | a00400ff             |

# MIPS Microarchitecture

## Multicycle μ architecture from Patterson & Hennessy



# Fibonacci (Binary)

- ◆ 1<sup>st</sup> statement: addi \$3, \$0, 8
- ◆ How do we translate this to machine language?
  - Hint: use instruction encodings below

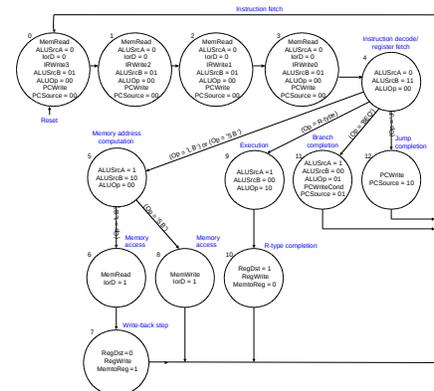
| format | example              | encoding   |     |    |       |      |    |    |    |     |    |    |   |       |
|--------|----------------------|--|-----|----|-------|------|----|----|----|-----|----|----|---|-------|
| R      | add \$rd, \$ra, \$rb | <table border="1"> <tr> <td>6</td> <td>5</td> <td>5</td> <td>5</td> <td>5</td> <td>6</td> </tr> <tr> <td>0</td> <td>ra</td> <td>rb</td> <td>rd</td> <td>0</td> <td>funct</td> </tr> </table> | 6   | 5  | 5     | 5    | 5  | 6  | 0  | ra  | rb | rd | 0 | funct |
| 6      | 5                    | 5  | 5   | 5  | 6     |      |    |    |    |     |    |    |   |       |
| 0      | ra                   | rb   | rd  | 0  | funct |      |    |    |    |     |    |    |   |       |
| I      | beq \$ra, \$rb, imm  | <table border="1"> <tr> <td>6</td> <td>5</td> <td>5</td> <td>16</td> </tr> <tr> <td>op</td> <td>ra</td> <td>rb</td> <td>imm</td> </tr> </table>  | 6   | 5  | 5     | 16   | op | ra | rb | imm |    |    |   |       |
| 6      | 5                    | 5  | 16  |    |       |      |    |    |    |     |    |    |   |       |
| op     | ra                   | rb   | imm |    |       |      |    |    |    |     |    |    |   |       |
| J      | j dest               | <table border="1"> <tr> <td>6</td> <td>26</td> </tr> <tr> <td>op</td> <td>dest</td> </tr> </table>   | 6   | 26 | op    | dest |    |    |    |     |    |    |   |       |
| 6      | 26                   |  |     |    |       |      |    |    |    |     |    |    |   |       |
| op     | dest                 |  |     |    |       |      |    |    |    |     |    |    |   |       |

# Fibonacci (Binary)

| Instruction       | Binary Encoding                       | Hexadecimal Encoding |
|-------------------|---------------------------------------|----------------------|
| addi \$3, \$0, 8  | 001000 00000 00011 0000000000001000   | 20030008             |
| addi \$4, \$0, 1  | 001000 00000 00100 0000000000000001   | 20040001             |
| addi \$5, \$0, -1 | 001000 00000 00101 1111111111111111   | 2005ffff             |
| beq \$3, \$0, end | 000100 00011 00000 0000000000000101   | 10600005             |
| add \$4, \$4, \$5 | 000000 00100 00101 00100 00000 100000 | 00852020             |
| sub \$5, \$4, \$5 | 000000 00100 00101 00101 00000 100010 | 00852822             |
| addi \$3, \$3, -1 | 001000 00011 00011 1111111111111111   | 2063ffff             |
| j loop            | 000010 0000000000000000000000000011   | 08000003             |
| sb \$4, 255(\$0)  | 110000 00000 00100 0000000011111111   | a00400ff             |

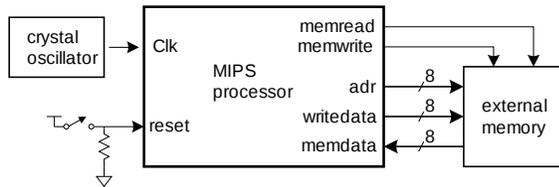
| format | example              | encoding  |     |    |       |      |    |    |    |     |    |   |       |
|--------|----------------------|---|-----|----|-------|------|----|----|----|-----|----|---|-------|
| R      | add \$rd, \$ra, \$rb | <table border="1"> <tr> <td>6</td> <td>5</td> <td>5</td> <td>5</td> <td>6</td> </tr> <tr> <td>0</td> <td>ra</td> <td>rb</td> <td>rd</td> <td>0</td> <td>funct</td> </tr> </table> | 6   | 5  | 5     | 5    | 6  | 0  | ra | rb  | rd | 0 | funct |
| 6      | 5                    | 5   | 5   | 6  |       |      |    |    |    |     |    |   |       |
| 0      | ra                   | rb  | rd  | 0  | funct |      |    |    |    |     |    |   |       |
| I      | beq \$ra, \$rb, imm  | <table border="1"> <tr> <td>6</td> <td>5</td> <td>5</td> <td>16</td> </tr> <tr> <td>op</td> <td>ra</td> <td>rb</td> <td>imm</td> </tr> </table>                                   | 6   | 5  | 5     | 16   | op | ra | rb | imm |    |   |       |
| 6      | 5                    | 5   | 16  |    |       |      |    |    |    |     |    |   |       |
| op     | ra                   | rb  | imm |    |       |      |    |    |    |     |    |   |       |
| J      | j dest               | <table border="1"> <tr> <td>6</td> <td>26</td> </tr> <tr> <td>op</td> <td>dest</td> </tr> </table>  | 6   | 26 | op    | dest |    |    |    |     |    |   |       |
| 6      | 26                   |   |     |    |       |      |    |    |    |     |    |   |       |
| op     | dest                 |   |     |    |       |      |    |    |    |     |    |   |       |

# Multicycle Controller

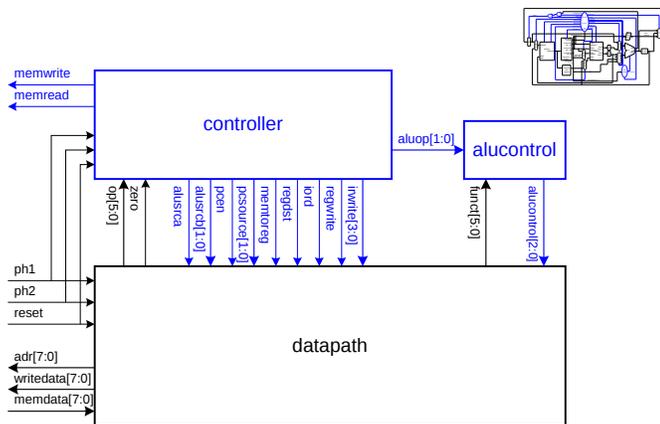


# Logic Design

- Start at top level
  - Hierarchically decompose MIPS into units
- Top-level interface



# Block Diagram



```

module controller(input clk, reset,
    input [5:0] op,
    input zero,
    output reg memread, memwrite, alusrcb,
    memtoreg, iord,
    output pcen,
    output reg regwrite, regdst,
    output reg [1:0] pcsource, alusrcb, aluop,
    output reg [3:0] irwrite);

parameter FETCH1 = 4'b0001;
parameter FETCH2 = 4'b0010;
parameter FETCH3 = 4'b0011;
parameter FETCH4 = 4'b0100;
parameter DECODE = 4'b0101;
parameter MEMADR = 4'b0110;
parameter LBRD = 4'b0111;
parameter LBWR = 4'b1000;
parameter SBWR = 4'b1001;
parameter RTYPEEX = 4'b1010;
parameter RTYPEWR = 4'b1011;
parameter BEQEX = 4'b1100;
parameter JEX = 4'b1101;
parameter ADDIWR = 4'b1110; // added for ADDI

parameter LB = 6'b100000;
parameter SB = 6'b101000;
parameter RTYPE = 6'b0;
parameter BEQ = 6'b000100;
parameter J = 6'b000010;
parameter ADDI = 6'b001000; // added for ADDI

reg [3:0] state, nextstate;
reg pcwrite, pcwritecond;
    
```

State Encodings...

## Controller Parameters

Opcodes...

Local reg variables...

# Verilog Code

```

// top level design includes both mips processor and memory
module mips_mem #(parameter WIDTH = 8, REGBITS = 3)(clk, reset);
    input clk, reset;
    wire memread, memwrite;
    wire [WIDTH-1:0] adr, writedata;
    wire [WIDTH-1:0] memdata;

    // instantiate the mips processor
    mips #(WIDTH,REGBITS) mips(clk, reset, memdata, memread,
        memwrite, adr, writedata);

    // instantiate memory for code and data
    exmem #(WIDTH) exmem(clk, memwrite, adr, writedata, memdata);
endmodule
    
```

```

// simplified MIPS processor
module mips #(parameter WIDTH = 8, REGBITS = 3)
    (input clk, reset,
    input [WIDTH-1:0] memdata,
    output memread, memwrite,
    output [WIDTH-1:0] adr, writedata);
    
```

## Top-level code

```

wire [31:0] instr;
wire zero, alusrcb, memtoreg, iord, pcen, regwrite, regdst;
wire [1:0] aluop, pcsource, alusrcb;
wire [3:0] irwrite;
wire [2:0] alucont;
    
```

```

controller cont(clk, reset, instr[31:26], zero, memread,
    memwrite,
    alusrcb, memtoreg, iord, pcen, regwrite, regdst,
    pcsource, alusrcb, aluop, irwrite);
alucontrol ac(aluop, instr[5:0], alucont);
datapath dp(WIDTH, REGBITS)
    dp(clk, reset, memdata, alusrcb, memtoreg, iord, pcen,
    regwrite, regdst, pcsource, alusrcb, irwrite, alucont,
    zero, instr, adr, writedata);
endmodule
    
```

# Main state machine – NS logic

```

// state register
always @(posedge clk)
    if(reset) state <= FETCH1;
    else state <= nextstate;

// next state logic (combinational)
always @(*)
    begin
        case(state)
            FETCH1: nextstate <= FETCH2;
            FETCH2: nextstate <= FETCH3;
            FETCH3: nextstate <= FETCH4;
            FETCH4: nextstate <= DECODE;
            DECODE: case(op)
                LB: nextstate <= MEMADR;
                SB: nextstate <= MEMADR;
                ADDI: nextstate <= MEMADR;
                RTYPE: nextstate <= RTYPEEX;
                BEQ: nextstate <= BEQEX;
                J: nextstate <= JEX;
                // should never happen
                default: nextstate <= FETCH1;
            endcase
        endcase

        MEMADR: case(op)
            LB: nextstate <= LBRD;
            SB: nextstate <= SBWR;
            ADDI: nextstate <= ADDIWR;
            // should never happen
            default: nextstate <= FETCH1;
        endcase

        LBRD: nextstate <= LBWR;
        LBWR: nextstate <= FETCH1;
        SBWR: nextstate <= FETCH1;
        RTYPEEX: nextstate <= RTYPEWR;
        RTYPEWR: nextstate <= FETCH1;
        BEQEX: nextstate <= FETCH1;
        JEX: nextstate <= FETCH1;
        ADDIWR: nextstate <= FETCH1;
        // should never happen
        default: nextstate <= FETCH1;
    end
endcase
    
```

# Setting Control Signal Outputs

```

always @(*)
begin
    // set all outputs to zero, then
    // conditionally assert just the
    // appropriate ones
    irwrite <= 4'b0000;
    pcwrite <= 0; pcwritecond <= 0;
    regwrite <= 0; regdst <= 0;
    memread <= 0; memwrite <= 0;
    alusrc_a <= 0; alusrc_b <= 2'b00;
    aluop <= 2'b00; pcsource <= 2'b00
    iord <= 0; memtoreg <= 0;
    case(state)
    FETCH1:
    begin
        memread <= 1;
        irwrite <= 4'b0001;
        alusrc_b <= 2'b01;
        pcwrite <= 1;
    end
    FETCH2:
    begin
        memread <= 1;
        irwrite <= 4'b0010;
        alusrc_b <= 2'b01;
        pcwrite <= 1;
    end
    FETCH3:
    begin
        memread <= 1;
        irwrite <= 4'b0100;
        alusrc_b <= 2'b01;
        pcwrite <= 1;
    end
    FETCH4:
    begin
        memread <= 1;
        irwrite <= 4'b1000;
        alusrc_b <= 2'b01;
        pcwrite <= 1;
    end
    DECODE: alusrc_b <= 2'b11;
    .....
    endcase
end
endmodule

```

# Verilog: alucontrol

```

module alucontrol(input [1:0] aluop,
input [5:0] funct,
output reg [2:0] alucont);

always @(*)
case(aluop)
2'b00: alucont <= 3'b010; // add for lb/sb/addi
2'b01: alucont <= 3'b110; // sub (for beq)
default: case(funct) // R-Type instructions
6'b100000: alucont <= 3'b010; // add (for add)
6'b100010: alucont <= 3'b110; // subtract (for sub)
6'b100100: alucont <= 3'b000; // logical and (for and)
6'b100101: alucont <= 3'b001; // logical or (for or)
6'b101010: alucont <= 3'b111; // set on less (for slt)
default: alucont <= 3'b101; // should never happen
endcase
endcase
endmodule

```

```

module datapath #(parameter WIDTH = 8, REGBITS = 3)
(input clk, reset,
input [WIDTH-1:0] memdata,
input alusrc_a, memtoreg, iord, pcen, regwrite, regdst,
input [1:0] pcsource, alusrc_b,
input [3:0] irwrite,
input [2:0] alucont,
output zero,
output [31:0] instr,
output [WIDTH-1:0] adr, writedata);

```

// the size of the parameters must be changed to match the WIDTH parameter

```

localparam CONST_ZERO = 8'b0;
localparam CONST_ONE = 8'b1;

```

```

wire [REGBITS-1:0] ra1, ra2, wa;
wire [WIDTH-1:0] pc, nextpc, md, rd1, rd2, wd, a, src1, src2, aluresult,
aluout, constx4;

```

```

// shift left constant field by 2
assign constx4 = {instr[WIDTH-3:0],2'b00};

```

```

// register file address fields
assign ra1 = instr[REGBITS+20:21];
assign ra2 = instr[REGBITS+15:16];
mux2 #(REGBITS) regmux(instr[REGBITS+15:16],
instr[REGBITS+10:11], regdst, wa);

```

## Verilog: Datapath 1

# Verilog: alu

```

module alu #(parameter WIDTH = 8)
(input [WIDTH-1:0] a, b,
input [2:0] alucont,
output reg [WIDTH-1:0] result);
wire [WIDTH-1:0] b2, sum, slt;

```

```

assign b2 = alucont[2] ? ~b:b;
assign sum = a + b2 + alucont[2];
// slt should be 1 if most significant bit of sum is 1
assign slt = sum[WIDTH-1];

```

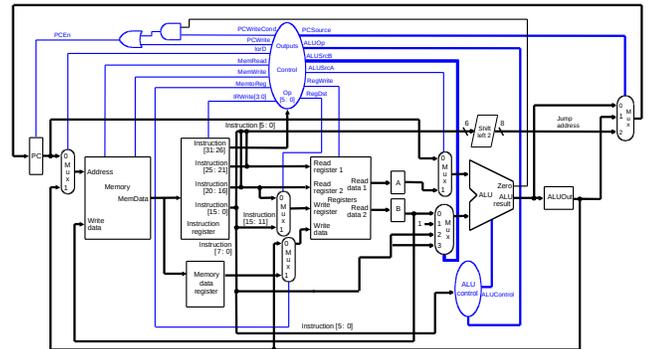
```

always@(*)
case(alucont[1:0])
2'b00: result <= a & b;
2'b01: result <= a | b;
2'b10: result <= sum;
2'b11: result <= slt;
endcase
endmodule

```

# MIPS Microarchitecture

- ◆ Multicycle  $\mu$  architecture from Patterson & Hennessy



// independent of bit width, load instruction into four 8-bit registers over four cycles

```

flop_en #(8) ir0(clk, irwrite[0], memdata[7:0], instr[7:0]);
flop_en #(8) ir1(clk, irwrite[1], memdata[7:0], instr[15:8]);
flop_en #(8) ir2(clk, irwrite[2], memdata[7:0], instr[23:16]);
flop_en #(8) ir3(clk, irwrite[3], memdata[7:0], instr[31:24]);

```

// datapath

```

flop_enr #(WIDTH) pcreg(clk, reset, pcen, nextpc, pc);
flop #(WIDTH) mdr(clk, memdata, md);
flop #(WIDTH) areg(clk, rd1, a);
flop #(WIDTH) wrd(clk, rd2, writedata);
flop #(WIDTH) res(clk, aluresult, aluout);
mux2 #(WIDTH) adrmux(pc, aluout, iord, adr);
mux2 #(WIDTH) src1mux(pc, a, alusrc_a, src1);
mux4 #(WIDTH) src2mux(writedata, CONST_ONE, instr[WIDTH-1:0],
constx4, alusrc_b, src2);
mux4 #(WIDTH) pcmux(aluresult, aluout, constx4, CONST_ZERO,
pcsource, nextpc);
mux2 #(WIDTH) wdmux(aluout, md, memtoreg, wd);
regfile #(WIDTH,REGBITS) rf(clk, regwrite, ra1, ra2, wa, wd, rd1, rd2);
alu #(WIDTH) alunit(src1, src2, alucont, aluresult);
zerodetect #(WIDTH) zd(aluresult, zero);
endmodule

```

## Verilog: Datapath 2

## Verilog: regfile

```

module regfile #(parameter WIDTH = 8, REGBITS = 3)
  (input      clk,
   input      regwrite,
   input [REGBITS-1:0] ra1, ra2, wa,
   input [WIDTH-1:0] wd,
   output [WIDTH-1:0] rd1, rd2);

  reg [WIDTH-1:0] RAM [(1<<REGBITS)-1:0];

  // three ported register file
  // read two ports (combinational)
  // write third port on rising edge of clock
  // register 0 is hardwired to 0
  always @(posedge clk)
    if (regwrite) RAM[wa] <= wd;

  assign rd1 = ra1 ? RAM[ra1] : 0;
  assign rd2 = ra2 ? RAM[ra2] : 0;
endmodule

```

## Verlog: Other stuff

```

module zerodetect #(parameter WIDTH
= 8)
  (input [WIDTH-1:0] a,
   output y);
  assign y = (a==0);
endmodule

module flop #(parameter WIDTH = 8)
  (input      clk,
   input      [WIDTH-1:0] d,
   output reg [WIDTH-1:0] q);
  always @(posedge clk)
    q <= d;
endmodule

module flopen #(parameter WIDTH = 8)
  (input      clk, en,
   input      [WIDTH-1:0] d,
   output reg [WIDTH-1:0] q);
  always @(posedge clk)
    if (en) q <= d;
endmodule

module flopenr #(parameter WIDTH = 8)
  (input      clk, reset, en,
   input      [WIDTH-1:0] d,
   output reg [WIDTH-1:0] q);
  always @(posedge clk)
    if (reset) q <= 0;
    else if (en) q <= d;
endmodule

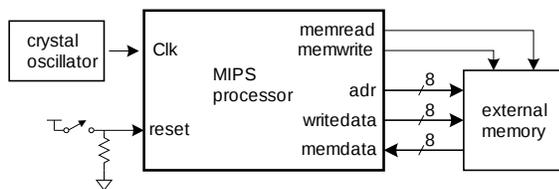
module mux2 #(parameter WIDTH = 8)
  (input [WIDTH-1:0] d0, d1,
   input      s,
   output [WIDTH-1:0] y);
  assign y = s ? d1 : d0;
endmodule

module mux4 #(parameter WIDTH = 8)
  (input [WIDTH-1:0] d0, d1, d2,
   d3,
   input [1:0] s,
   output reg [WIDTH-1:0] y);
  always @(*)
  case(s)
    2'b00: y <= d0;
    2'b01: y <= d1;
    2'b10: y <= d2;
    2'b11: y <= d3;
  endcase
endmodule

```

## Logic Design

- Start at top level
  - Hierarchically decompose MIPS into units
- Top-level interface



## Verilog: exmemory

```

// external memory accessed by MIPS
module exmemory #(parameter
WIDTH = 8)
  (clk, memwrite, adr, writedata,
   memdata);

  input      clk;
  input      memwrite;
  input [WIDTH-1:0] adr,
  writedata;
  output reg [WIDTH-1:0] memdata;

  reg [31:0] RAM [(1<<WIDTH-2)-
1:0];
  wire [31:0] word;

  initial
  begin

    $readmemh("memfile.dat",RAM);
  end

  // read and write bytes from 32-bit word
  always @(posedge clk)
  if(memwrite)
  case (adr[1:0])
    2'b00: RAM[adr>>2][7:0] <=
writedata;
    2'b01: RAM[adr>>2][15:8] <=
writedata;
    2'b10: RAM[adr>>2][23:16] <=
writedata;
    2'b11: RAM[adr>>2][31:24] <=
writedata;
  endcase

  assign word = RAM[adr>>2];
  always @(*)
  case (adr[1:0])
    2'b00: memdata <= word[7:0];
    2'b01: memdata <= word[15:8];
    2'b10: memdata <= word[23:16];
    2'b11: memdata <= word[31:24];
  endcase
endmodule

```

## Synthesized memory?

- If you synthesize the Verilog, you'll get a memory
  - But – it will be huge!
  - It will be made of your DFF cells
  - plus synthesized address decoders
  - Custom memory is much smaller
  - but much trickier to get right
  - ... see details in VGA slides ...

## Verilog: exmemory

```

// external memory accessed by
MIPS
module exmem #(parameter WIDTH
= 8)
  (clk, memwrite, adr,
   writedata, memdata);

  input      clk;
  input      memwrite;
  input [WIDTH-1:0] adr,
  writedata;
  output [WIDTH-1:0] memdata;

  wire memwriteB, clkB;

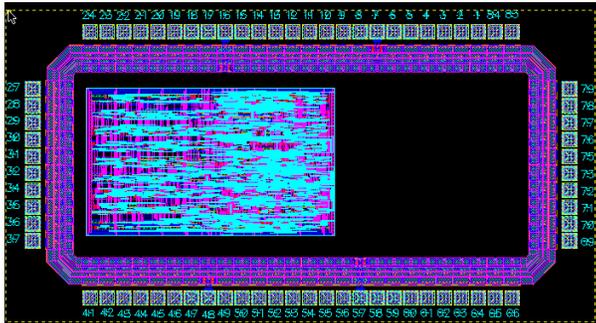
  // UMC RAM has active low write
  enable...
  not(memwriteB, memwrite);

  // Looks like you need to clock the
  memory early
  // to make it work with the current
  control...
  not(clkB, clk);

  // Instantiate the UMC SPRAM module
  UMC130SPRAM_8_8_mips_ram (
    .CK(clkB),
    .CEN(1'b0),
    .WEN(memwriteB),
    .OEN(1'b0),
    .ADR(adr),
    .DI(writedata),
    .DOUT(memdata));
endmodule

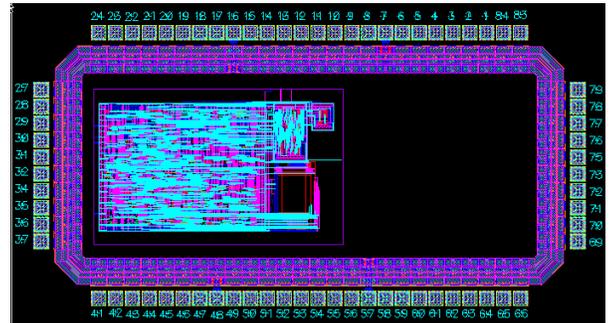
```

## MIPS (8-bit) size comparison



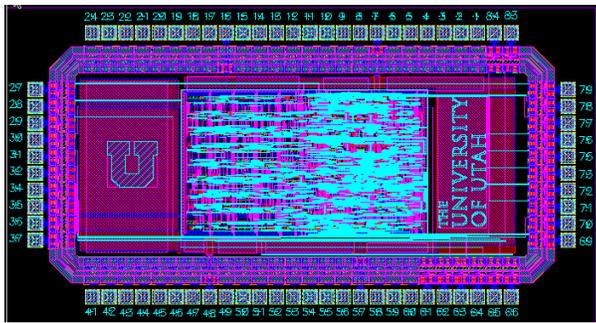
One big EDI run of the whole thing  
With some work, could probably get this in a single TCU...

## MIPS (8-bit) size comparison



Separate EDI for controller, alucontrol, datapath  
and custom RF, assembled with ccar

## MIPS (8-bit) whole chip



Includes poly/m1/m2/m3 fill, and logos  
Routed to the pads using ccar